

Lecture 2

Algorithmic Number Theory for Function Fields

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First Part

Consider complete regular curves C over a field K . We can then equivalently work with $F = K(C)$ only.

Notation:

- ▶ Group of divisors $\text{Div}(F/K)$.
- ▶ Subset of divisors of degree d : $\text{Div}^d(F/K)$.
- ▶ Subgroup of principal divisors of $\text{Princ}(F/K)$.
- ▶ Class group or Picard group: $\text{Pic}(F/K)$.
- ▶ Subgroup of class of degree d : $\text{Pic}^d(F/K)$.

By definition,

$$\text{Pic}(F/K) = \text{Div}(F/K)/\text{Princ}(F/K),$$

$$\text{Pic}^0(F/K) = \text{Div}^0(F/K)/\text{Princ}(F/K).$$

We have

$$\text{Pic}(F/K) = \text{Pic}^0(F/K) \oplus \langle A \rangle,$$

where A is a divisor of F/K with minimal positive degree.

If K is a finite field or algebraically closed then $\deg(A) = 1$. In the latter case A can be chosen to be a prime divisor.

If K is finitely generated over its prime field then $\text{Pic}^0(F/K)$ is finitely generated.

If K is a finite field then $\text{Pic}^0(F/K)$ is finite. Then usually

$$\#\text{Pic}^0(F/K) \approx (\#K)^g.$$

Computing in the Class Group

Representation of divisors:

- ▶ Divisors can be represented as a sum of places with integral coefficients, or as a pair of fractional ideals.
- ▶ Addition of divisors either by addition of coefficient vectors or multiplication of ideals.
- ▶ Equality by coefficientwise comparison or comparison of Hermite normal forms.

Representation of divisor classes:

- ▶ By divisors, which can be “suitably” chosen, for example reduced divisors.
- ▶ Comparison via unique divisor class representatives, if they can be computed, or by the test

$$\deg(D) = \deg(E) \text{ and } L(D - E) \neq 0.$$

- ▶ This is usually efficient (polynomial time) in terms of operations in K .

Reduction of divisors:

- ▶ Fix a divisor A of positive degree.
- ▶ For every D there is $\tilde{D} \geq 0$ and $f \in F^\times$ such that

$$D = \tilde{D} - rA + \operatorname{div}(f)$$

and $\deg(\tilde{D}) \leq g + \deg(A) - 1$.

- ▶ If A is a prime divisor of degree one and r is minimal then \tilde{D} is uniquely determined.

Class representatives:

- ▶ Thus $[D] = [\tilde{D} - rA]$ for every divisor class.
- ▶ If A is a prime divisor of degree one then $\tilde{D} - rA$ can be uniquely chosen.

Computing in the Class Group

Reduction of divisors:

- ▶ The reduced divisor \tilde{D} can be computed by an iterative double-and-add method such that the runtime is polynomial in g , $\deg(A)$ and the length of D .
- ▶ Moreover, f is computed as a product of powers of elements of F and has length polynomial in g , $\deg(A)$ and the length of D .

These ideas can be optimised, for example by precomputations, and worked out in great detail.

A (biased) selection of results:

- ▶ Cantor: If F/K is hyperelliptic then operations in $\text{Pic}^0(F/K)$ can be reduced to fast polynomial arithmetic in degree $O(g)$, so the runtime is $O^\sim(g)$.
- ▶ Makdisi: If F/K is arbitrary then operations in $\text{Pic}^0(F/K)$ can be reduced to fast matrix arithmetic in dimension $O(g)$, so the runtime is $O^\sim(g^\omega)$.
- ▶ Hess-Junge: If F/K has a rational subfield of index n , where $n = O(g)$ is always possible, then operations in $\text{Pic}^0(F/K)$ can be reduced to fast polynomial matrix arithmetic in dimension $O(n)$ and degree $O(g/n)$, so the runtime is $O^\sim(n^\omega(g/n))$.

Computing the Class Group

We assume that K is finite! Write $q = \#K$.

Have $\text{Pic}^0(F/K) \cong \mathbb{Z}/c_1\mathbb{Z} \times \cdots \times \mathbb{Z}/c_{2g}\mathbb{Z}$. with $c_i | c_{i+1}$.

Goal:

- ▶ Compute the c_i .
- ▶ Compute images and preimages under a fixed isomorphism

$$\phi : \text{Pic}(F/K) \rightarrow \mathbb{Z} \oplus \mathbb{Z}/c_1\mathbb{Z} \times \cdots \times \mathbb{Z}/c_{2g}\mathbb{Z}.$$

Denote by A a fixed divisor of degree one that maps under ϕ to the first cyclic factor of the codomain of ϕ .

Computing the Class Group

Algorithms that work for any finite abelian group G :

- ▶ Classic runtime $O((\#G)^{1/2})$.
- ▶ Improvements often lead to $O((\#G)^{1/3})$.
- ▶ So here roughly $O^{\sim}(q^{g/2})$ or $O^{\sim}(q^{g/3})$.

Algorithms that use $G = \text{Pic}^0(F/K)$ usually employ an index calculus strategy:

- ▶ If q is small and g is large, the (heuristic) runtime is $q^{(c+o(1))g^{1/2} \log(g)^{1/2}}$, and $q^{(d+o(1))g^{1/3} \log(g)^{2/3}}$ (†) in special families.
- ▶ If q is large and $g \geq 2$ fixed, then $O^{\sim}(q^{2-2/g})$ (†).

(†): This is for discrete logarithms, so restrictions may apply.

Index Calculus

Setup:

- ▶ Let S denote the set of places of F/K of degree $\leq r$, called factor basis.
- ▶ Let $[D_1], \dots, [D_s]$ denote generators of $\text{Pic}^0(F/K)$.

Relation search:

- ▶ Choose random λ_i and compute $[\tilde{D} - IA] = \sum_i \lambda_i [D_i]$ with \tilde{D} reduced.
- ▶ Factor \tilde{D} over S , if possible and obtain

$$\sum_i \lambda_i [D_i] = [\tilde{D} - IA] = -I[A] + \sum_{P \in S} n_P [P].$$

- ▶ Store λ_i and n_P as rows of a matrix and repeat.

Index Calculus

Linear algebra:

- ▶ If matrix has full rank and sufficiently more rows than columns use a Hermite normal form computation to derive relations between the generators $[D_i]$.
- ▶ Use a Smith normal form computation to derive c_1, \dots, c_{2g} from those relations.

Why does it work?

- ▶ There is a good upper bound on r .
- ▶ The class number can be efficiently approximated and checked against the computed c_1, \dots, c_{2g} .
- ▶ There is a reasonable good (heuristic) probability that enough relations are obtained.

Index Calculus

Let N_m denote the number of places of degree one in the constant field extension of F of degree m .

Theorem. Suppose $N_r > (g - 1)2q^{r/2}$. Then $\text{Pic}(F/K)$ is generated by the places of degree $\leq r$ and the places in $\text{supp}(A)$.

Theorem. Let $h = \#\text{Pic}^0(F/K)$. Then

$$\left| \log\left(\frac{h}{q^g}\right) - \sum_{m=1}^t \frac{q^{-m}}{m} (N_m - q^m - 1) \right| \leq \frac{2g}{q^{1/2} - 1} \cdot \frac{q^{-t/2}}{t + 1}.$$

Since $c_1 \dots c_{2g}$ is an integral multiple of h an approximation of $\log(h/q^g)$ up to an error of $\log(2)/3$ is sufficient.

Some Applications

From the relations of the $[D_i]$ it is easy to compute generators of $\text{Pic}^0(F/K)$ corresponding to the cyclic generators of the codomain of ϕ . We can thus also compute preimages under ϕ efficiently.

Images $\phi([D])$ are computed by adding $[D - \deg(D)A]$ to the $[D_i]$ and searching for relations. The runtime is then basically the same like that for computing the c_1, \dots, c_{2g} .

This can directly be used to compute for an arbitrary S

- ▶ S -units $U(S) = \{f \in F^\times \mid \text{supp}(\text{div}(f)) \subseteq S\}$ and
- ▶ S -class groups $\text{Div}(F/K)/(\langle S \rangle + \text{Princ}(F/K))$.

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Second Part

Notation:

- ▶ Let \mathfrak{m} denote an effective divisor, called modulus.
- ▶ $\text{Div}_{\mathfrak{m}}(F/K)$ group of divisors coprime to \mathfrak{m} .
- ▶ $F_{\mathfrak{m}}^{\times} = \{f \in F^{\times} \mid v_P(f - 1) \geq v_P(\mathfrak{m}) \text{ for all } P\}$ group of elements congruent to one modulo \mathfrak{m} .
- ▶ $\text{Princ}_{\mathfrak{m}}(F/K) = \{\text{div}(f) \mid f \in F_{\mathfrak{m}}^{\times}\}$, the ray modulo \mathfrak{m} .
- ▶ $\text{Pic}_{\mathfrak{m}}(F/K) = \text{Div}_{\mathfrak{m}}(F/K) / \text{Princ}_{\mathfrak{m}}(F/K)$, the ray class group modulo \mathfrak{m} .
- ▶ $\phi_{\mathfrak{m},n} : \text{Pic}_{\mathfrak{m}}(F/K) \rightarrow \text{Pic}_n(F/K)$, $[D]_{\mathfrak{m}} \mapsto [D]_n$ for $n \geq \mathfrak{m}$.

We have $\text{Princ}_{\gcd(\mathfrak{m},n)}(F/K) = \text{Princ}_{\mathfrak{m}}(F/K) + \text{Princ}_n(F/K)$.

The $\phi_{\mathfrak{m},n}$ are epimorphisms.

Artin Map

Let E/F be a finite abelian extension. Let P be place of F/K and write $N(P) = \#\mathcal{O}_P/\mathfrak{m}_P = q^{\deg(P)}$.

If P is unramified in E/F then there is a uniquely determined $\sigma_P \in \text{Gal}(E/F)$ satisfying

$$\sigma_P(x) \equiv x^{N(P)} \pmod{\mathfrak{m}_Q}$$

for all places Q of E/K above P and all $x \in \mathcal{O}_Q$.

Suppose E/F is unramified outside $\text{supp}(\mathfrak{m})$. The Artin map is defined as

$$A_{E/F} : \text{Div}_{\mathfrak{m}}(F/K) \rightarrow \text{Gal}(E/F), \quad D \mapsto \prod_P \sigma_P^{v_P(D)}.$$

Some Properties of the Artin Map

Theorem.

- ▶ The Artin map is surjective.
- ▶ If the multiplicities of \mathfrak{m} are large enough then

$$\text{Princ}_{\mathfrak{m}}(F/K) \subseteq \ker(A_{E/F}).$$

Any \mathfrak{m} like in the theorem is called a modulus of E/F . There is a smallest modulus $\mathfrak{f}(E/F)$ of E/F , called conductor of E/F . Every place in \mathfrak{m} is ramified in E/F .

If \mathfrak{m} is a modulus of E/F then regard

$$A_{E/F} : \text{Pic}_{\mathfrak{m}}(F/K) \rightarrow \text{Gal}(E/F).$$

Thus if $H = \ker(A_{E/F})$ then H has finite index in $\text{Pic}_{\mathfrak{m}}(F/K)$ and

$$\text{Gal}(E/F) \cong \text{Pic}_{\mathfrak{m}}(F/K)/H.$$

Define

$$N_{E/F} : \text{Pic}_{\text{Con}_{E/F}(\mathfrak{m})}(E/K) \rightarrow \text{Pic}_{\mathfrak{m}}(F/K)$$

by taking the norm of a representing divisor. Norms of elements of $E_{\text{Con}_{E/F}(\mathfrak{m})}^{\times}$ are elements of $F_{\mathfrak{m}}^{\times}$, so this is well defined.

Theorem. If E/F is finite abelian with modulus \mathfrak{m} then

$$\ker(A_{E/F}) = \text{im}(N_{E/F}).$$

We say that E is a class field over F with modulus \mathfrak{m} that belongs to the subgroup $H = \text{im}(N_{E/F}) = \ker(A_{E/F})$ of finite index of $\text{Pic}_{\mathfrak{m}}(F/K)$.

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Theorem.

1. If H is any subgroup of $\text{Pic}_{\mathfrak{m}}(F/K)$ of finite index, then there is a class field E over F with modulus \mathfrak{m} that belongs to H , and E is uniquely determined up to F -isomorphism.
2. The degree of the exact constant field of E/K over K is equal to $\deg(H)$, the minimal positive degree of divisor classes in H .

Computing Ray Class Groups

There is an exact sequence of finitely generated abelian groups

$$0 \rightarrow K^\times \rightarrow \prod_P (\mathcal{O}_P / \mathfrak{m}_P^{v_P(\mathfrak{m})})^\times \rightarrow \text{Pic}_m(F/K) \rightarrow \text{Pic}(F/K) \rightarrow 0.$$

We have:

- ▶ Generators and relations can be computed for each object of the sequence other $\text{Pic}_m(F/K)$.
- ▶ Elements of each object can be represented in chosen generators.
- ▶ Images and preimages of the maps of the sequence can also be computed.

Then generators and relations of $\text{Pic}_m(F/K)$ can be computed and elements of $\text{Pic}_m(F/K)$ can be represented in those generators.

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Given $H \leq \text{Pic}_m(F/K)$ the goal is to compute defining equations for the class field E over F of modulus m that belongs to H .

Theorem. Suppose $H_1, H_2 \subseteq \text{Pic}_m(F/K)$ with $H_1 \cap H_2 = H$. If E_1 belongs to H_1 and E_2 belongs to H_2 then $E = E_1 E_2$ belongs to H .

We can choose H_1 and H_2 such that the index of H_1 is coprime to $\text{char}(F)$ and the index of H_2 is a power of $\text{char}(F)$.

Coprime to Characteristic Case

Theorem. Let F'/F finite and $E' = EF'$. Then E' is the class field over F' with modulus $\mathfrak{m}' = \text{Con}_{F'/F}(\mathfrak{m})$ that belongs to $H' = N_{F'/F}^{-1}(H)$.

Suppose that the index of H is coprime to $\text{char}(F)$ and let n denote the exponent of $\text{Pic}_{\mathfrak{m}}(F/K)/H$.

Let $F' = F(\mu_n)$.

Theorem. Every abelian extension of F' of exponent n is a Kummer extension, is thus obtained by adjoining n -th roots of suitable Kummer elements of F' to F' .

This leads to a rather explicit representation of E' .

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Then it is known and can be done:

- ▶ Kummer elements f_i can be computed for the class field G over F' of modulus \mathfrak{m}' that belongs to $n\text{Pic}_{\mathfrak{m}}(F'/K)$, for example by an S -units computation in F' .
- ▶ H' is computed as a preimage of maps of abelian groups.
- ▶ E' is the fixed field of G under $A_{G/F'}(H')$, the Kummer elements g_j of E' are accordingly computed as products of the f_i using a generalised Tate-Lichtenbaum pairing.
- ▶ E'/F is finite abelian with modulus m , and E is the fixed field of E' under $A_{E'/F}(H)$. Defining equations for E can be computed via explicit Galois theory.

Power of Characteristic Case

Theorem. Every abelian extension of F' of exponent n , an m -th power of $\text{char}(F)$, is an Artin-Schreier-Witt extension, is thus obtained by adjoining the division points of A-S-W elements in $W_m(F')$ under the A-S-W operator to F' .

This leads to a rather explicit but also rather involved representation of E . Let n be the exponent of $\text{Pic}_m(F/K)/H$.

Then it is known and can be done:

- ▶ A-S-W elements f_i can be computed for the class field G over F of modulus \mathfrak{m} that belongs to $n\text{Pic}_m(F/K)$, for example by a Riemann-Roch computation in F .
- ▶ E is the fixed field of G under $A_{G/F}(H)$, the A-S-W elements g_j of E are accordingly computed as sums of the f_i using a pairing.

Applications

Construction of function fields with many rational places:

- ▶ A place P of F is fully split in E if and only if $P \in H$.
- ▶ Let $h_{n,H} = \#\text{Pic}_n(F/K)/\phi_{m,n}(H)$. The genus of E satisfies

$$\deg(H)(g_E - 1) = h_{m,H} \left(g_F - 1 + \frac{\deg(\mathfrak{m})}{2} \right) - \frac{1}{2} \sum_{P|\mathfrak{m}} \left(\sum_{k=1}^{v_P(\mathfrak{m})} h_{m-kP,H} \right) \deg(P).$$

Construction of Drinfeld modules:

- ▶ Is defined by coefficients which are elements of a specific class field.
- ▶ The coefficients satisfy various relations.
- ▶ Use those relations to solve for the coefficients over the class field.

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Third Part

Motivation

Study zero sets of polynomial equations over various fields

- ▶ Example: $\{(x, y) \in K^2 \mid x^2 + y^2 = 1\}$
- ▶ Over finite fields: Count solutions!

Algebraic curves: Polynomial equations have one free variable, the other variables are algebraically dependent.

We will again consider function fields F/\mathbb{F}_q over the exact constant field \mathbb{F}_q instead of curves. Write N_d for the places of degree one of F/\mathbb{F}_q .

The zeta function of F/K is

$$\begin{aligned} \zeta_{F/K}(t) &= \exp\left(\sum_{d=1}^{\infty} N_d \cdot \frac{t^d}{d}\right) \\ &= \prod_P \frac{1}{1 - t^{\deg(P)}} = \sum_{D \geq 0} t^{\deg(D)}. \end{aligned}$$

Frobenius Operation

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There is $L_{F/K}(t) \in \mathbb{Z}[t]$ with $\deg(L_{F/K}(t)) = 2g$ and

$$\zeta_{F/K}(t) = \frac{L_{F/K}(t)}{(1-t)(1-qt)}.$$

This is called the L -polynomial of F/K .

Moreover, there are \mathbb{Q}_ℓ -vector spaces V_ℓ and $\text{Frob}_{q,\ell} \in \text{Aut}(V_\ell)$ such that

$$L_{F/K}(t) = \det(\text{id} - \text{Frob}_{q,\ell} \cdot t \mid V_\ell).$$

Computation of Zeta functions

Possible applications:

- ▶ “Cryptography”
- ▶ Distribution of the eigenvalues of Frobenius
- ▶ ...

Complexity of ℓ -adic methods:

- ▶ Exponential in g and polynomial in $\log(q)$,
- ▶ impractical for $g \geq 3$.

Complexity of p -adic methods:

- ▶ Mostly $O^{\sim}(p^1 g^4 n^3)$ or $O^{\sim}(p^1 g^5 n^3)$ with $n = \log_p(q)$.
- ▶ Random $q = 2, g = 350$ hyperelliptic curve in 3 days.

Galois and Abelian Extensions

Let E/F denote a finite Galois extension with Galois group G such that K is the exact constant field of E .

The associated product formula for $\zeta_{E/K}(t)$ is

$$\zeta_{E/K}(t) = \prod_{\chi} L(E/F, \chi, t)^{\chi(1)},$$

where χ runs over the irreducible characters of G and $L(E/F, \chi, t)$ will be defined later (for G abelian).

Can the product be computed more efficiently for large g_E ?

If E/F is abelian then E is a class field over F belonging to some H and the factors of the product can be described in terms of H !

Ray Class Groups

We have already met ray class groups. Here are some (more) properties.

For a subgroup H of $\text{Pic}_{\mathfrak{m}}(F/K)$ of finite index there is a unique minimal $\mathfrak{f}(H) \leq \mathfrak{m}$ with

$$\text{Pic}_{\mathfrak{f}(H)}(F/K) / \phi_{\mathfrak{m}, \mathfrak{f}(H)}(H) \cong \text{Pic}_{\mathfrak{m}}(F/K) / H.$$

The divisor $\mathfrak{f}(H)$ is the conductor of H . It is equal to the conductor of the class field E over F belonging to H .

$$\text{Pic}_{\mathfrak{m}}(F/K) \cong \text{Pic}_{\mathfrak{m}}^0(F/K) \oplus \mathbb{Z}.$$

$$\#\text{Pic}_{\mathfrak{m}}^0(F/K) = \frac{\#\text{Pic}^0(F/K) \cdot \prod_{i=1}^s (q^{\deg(P)} - 1) q^{\deg(P)(v_P(\mathfrak{m})-1)}}{q - 1}.$$

Characters and L-series

A character χ modulo \mathfrak{m} is a homomorphism

$$\chi : \text{Pic}_{\mathfrak{m}}(F/K) \rightarrow \mathbb{C}^{\times}$$

of finite order. The conductor $f(\chi)$ of χ is $f(\ker(\chi))$.

The character sum $N_d(\chi)$ of degree d is

$$N_d(\chi) = \sum_{\deg(P)|d, P \not\leq f(\chi)} \deg(P) \cdot \chi([P])^{d/\deg(P)}.$$

The L-series $L(\chi, t) = L(E/F, \chi, t)$ of χ with $\ker(\chi) \supseteq H$ is

$$L(\chi, t) = \exp \left(\sum_{d=1}^{\infty} N_d(\chi) \cdot t^d / d \right).$$

We have $\zeta_{F/K}(t) = L(\chi, t)$ for $\chi = \text{id}$.

Theorem. Assume $\ker(\chi) \neq \text{Pic}_m(F/K)$. Then

$$L(\chi, t) = \prod_{i=1}^{2g-2+\deg(f(\chi))} (1 - \omega_i(\chi)t)$$

with $|\omega_i(\chi)| = q^{1/2}$ and ζ primitive $\text{ord}(\chi)$ -th root of unity, and

$$L(\chi, t) = \varepsilon(\chi) \cdot q^{g-1+\deg(f(\chi))/2} \cdot t^{2g-2+\deg(f(\chi))} \cdot L(\bar{\chi}, \frac{1}{qt})$$

with $\varepsilon(\chi) \in q^{-\deg(f(\chi))/2} \mathbb{Z}[\zeta]$ and $|\varepsilon(\chi)| = 1$. Furthermore,

$$\begin{aligned} \zeta_{E/K}(t) &= \frac{L_{E/K}(t)}{(1-t)(1-qt)} \\ &= \frac{L_{E/K}(t) \cdot \prod_{\text{Pic}_m(F/K) \not\supseteq \ker(\chi) \supseteq H} L(\chi, t)}{(1-t)(1-qt)} \end{aligned}$$

Computing one L-series

Let $L(\chi, t) = \sum_{i=0}^{2g-2+\deg f(\chi)} a_i t^i$ with $a_i \in \mathbb{Z}[\zeta]$ and $a_0 = 1$.

1. The coefficients a_1, \dots, a_m can be computed from $N_1(\chi), \dots, N_m(\chi)$ by the definition of $L(\chi, t)$:

$$L(\chi, t) = \sum_{i=0}^m a_i t^i \equiv \exp \left(\sum_{d=1}^m N_d(\chi) \cdot t^d / d \right) \pmod{t^{m+1}}.$$

2. The character sums $N_1(\chi), \dots, N_m(\chi)$ can be computed from their definition

$$N_d(\chi) = \sum_{\deg(P)|d, P \not\leq f(\chi)} \deg(P) \cdot \chi([P])^{d/\deg(P)}$$

by enumerating all places P up to degree m with $P \not\leq f(\chi)$.

Computing one L-series

3. Compute characters χ modulo \mathfrak{m} with $\ker(\chi) \supseteq H$:
 - ▶ Use representations of $\text{Pic}_{\mathfrak{m}}(F/K)$, H and $\text{Pic}_{\mathfrak{m}}(F/K)/H$ in terms of generators and relations.
 - ▶ Define χ on generators of $\text{Pic}_{\mathfrak{m}}(F/K)/H$ and pull back to $\text{Pic}_{\mathfrak{m}}(F/K)$.
 - ▶ Compute $\ker(\chi) \supseteq H$ and $f(\chi) = f(\ker(\chi))$.
 - ▶ Write P in the generators of $\text{Pic}_{f(\chi)}(F/K)$ to obtain $\chi([P])$.

4. Due to the functional equation there is some redundancy between the coefficients of $L(\chi, t)$. As a consequence it often suffices to take m about half the degree of $L(\chi, t)$.

Best to have a toolbox for finitely generated abelian groups and homomorphisms. Requires algorithms for structure computation of $\text{Pic}_{\mathfrak{m}}(F/K)$ and discrete logarithms in $\text{Pic}_{\mathfrak{m}}(F/K)$.

Computing the Zeta function

Need to choose one ζ for all χ on $\text{Pic}_m(C)$ with $\ker(\chi) \supseteq H$.

Compute $L_{E/K}(t)$ as product over all L -series

$$L_{E/K}(t) = L_{F/K}(t) \cdot \prod_{\text{Pic}_m(F/K) \supseteq \ker(\chi) \supseteq H} L(\chi, t).$$

Use some optimisations:

- ▶ Let $\sigma \in \text{Gal}(\mathbb{Q}(\zeta)/\mathbb{Q})$. Then $L(\sigma \circ \chi, t) = L(\chi, t)^\sigma$. Use Galois redundancy: Compute system of representatives R for $\text{Gal}(\mathbb{Q}(\zeta)/\mathbb{Q})$ -orbits of $(\text{Pic}_m(F/K)/H)^*$. For each $\chi \in R$ compute $L(\chi, t)$ and derive $L(\sigma \circ \chi, t) = L(\chi, t)^\sigma$.
- ▶ Choose some epimorphism $\psi : \mathbb{Z}[\zeta] \rightarrow \mathbb{Z}/n\mathbb{Z}$ with n large. Compute product over $\mathbb{Z}/n\mathbb{Z}$ and reconstruct coefficients of $L_{E/K}(t)$ from $\mathbb{Z}/n\mathbb{Z}$ to \mathbb{Z} by choosing the representative of smallest absolute value.

In the following only very rough estimations.

Input size: $F/K, m, H$ polynomial in $\log(q), g, \deg(m)$.

Output size: $g_E^2 \log(q)$.

Computing one L-series: $q^{2(g+\deg(f(x)))}$.

Computing Zeta function:

- ▶ L-series product: $g_E^2 \log(q)$.
- ▶ Galois redundancy gives big practical, but no asymptotic speed up.

Depending on H have very roughly $\deg(m) \lesssim g_E \lesssim q^{g+\deg(m)}$.

So for small H asymptotically optimal!

Galois module structure of $\text{Pic}^0(E/K)$:

- ▶ Use L -series to compute Stickelberger element in the group ring $\mathbb{Z}[G]$
- ▶ Derive information about the structure of $\text{Pic}^0(E/K)$ via Stickelberger ideal and Kolyvagin derivative classes.
- ▶ Derive relations of conjugate elements in $\text{Pic}^0(E/K)$ under certain conditions.

This is interesting since no equations for E/K and no expensive class group computation of $\text{Pic}^0(E/K)$ needs to be carried out.

1. Show that there is an injective map of sets of $\text{Pic}^0(F/K)$ into the set of effective divisors of degree n , for any $n \geq n$.

2. Show that $\text{Pic}^0(K(x)/K) = 0$.

3. Show that $\text{Pic}_m(F/K) \cong \text{Pic}(F/K)$ if and only if m is a prime divisor of degree one.

4. Let $\phi : E_1 \rightarrow E_2$ be a morphism of elliptic curves. Show that $K(E_1)$ is a class field of $\phi^*(K(E_2))$ belonging to

$$H = \langle \infty \rangle \times \{(\phi(P)) - (\infty) \mid P \in E_1(K)\}.$$

5. If $\chi \neq 1$ is a character for $\mathbb{F}_q(x)/\mathbb{F}_q$ then $\deg(f(\chi)) \geq 2$.

6. Let $F = \mathbb{F}_7(x, y)$ with $y^2 = x^5 + 2x + 1$. Compute the genus and number of rational places of the class field of F/K with modulus $\mathfrak{m} = 2\infty + 3(x, y - 1)$ and subgroup H generated by $[(x, y + 1)]_{\mathfrak{m}}$.